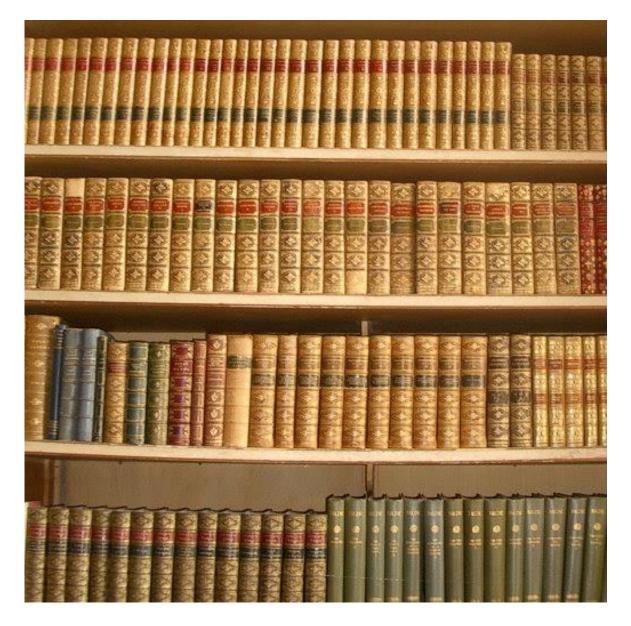
## Searching and sorting



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## Sequential search

- Sequential search
  - Scan through array, looking for key.
  - Search hit: return array index.
  - Search miss: return -1.

```
public static int search(String key, String[] a)
{
    for (int i = 0; i < a.length; i++)
        if (a[i].compareTo(key) == 0)
            return i;
    return -1;
}</pre>
```

# Search client, exception filter

- Exception filter
  - Read sorted list of strings from a whitelist file
  - Print strings from standard input not in whitelist

```
public static void main(String [] args)
{
    In in = new In(args[0]);
    String s = in.readAll();
    String[] words = s.split("\\s+");
    while (!StdIn.isEmpty())
    {
        String key = StdIn.readString();
        if (search(key, words) == -1)
            System.out.println(key);
    }
}
```

% more test.txt bob@office carl@beach marvin@spam bob@office bob@office mallory@spam dave@boat eve@airport alice@home

```
% more whitelist.txt
alice@home
bob@office
carl@beach
dave@boat
```

% java Whitelist whitelist.txt < test.txt
marvin@spam
mallory@spam
eve@airport

# Searching challenge 1

- Problem: A credit card company needs to whitelist 10 million customer account numbers, processing 10,000 transactions per second
- Question: Using <u>sequential search</u>, what kind of computer is needed?
  - A. Toaster.
  - B. Cell phone.
  - C. Your laptop.
  - D. Supercomputer. E. Google server farm.

# **Binary search**

#### Main idea 10 aback Sort the array (stay tuned) Play "20 questions" to determine index with a given key the key mid macabre (known value) Examples: Dictionary, phone book, is between book index, credit card numbers, a[mid] and a[hi-1] ? query the index ... (unknown value) is between mid and hi-1 hi-1 zygote

- Binary search
  - Examine the middle key.
  - If it matches, return its index.
  - Otherwise, search either the left or right half.

Binary search in a sorted array (one step)

# Binary search: Java implementation

#### Invariant

- Algorithm maintains:  $a[10] \leq key \leq a[hi-1]$ 

• Java library implementation: Arrays.binarySearch()

```
public static int search(String key, String[] a)
{
    return search(key, a, 0, a.length);
}
public static int search(String key, String[] a, int lo, int hi)
{
    if (hi <= lo) return -1;
    int mid = lo + (hi - lo) / 2;
    int cmp = a[mid].compareTo(key);
    if (cmp > 0) return search(key, a, lo, mid);
    else if (cmp < 0) return search(key, a, mid+1, hi);
    else return mid;
}</pre>
```

"I was amazed: given ample time, only about ten percent of professional programmers were able to get this small program right. But they aren't the only ones to find this task difficult: in the history in Section 6.2.1 of his Sorting and Searching, Knuth points out that while the first binary search was published in 1946, the first published binary search without bugs did not appear until 1962."

- Jon Bentley, Programming Pearls

## Binary search: mathematical analysis

- Analysis, binary search array of size N
  - Do one compare
  - Then binary search in an array of size N/2

$$- N \rightarrow N/2 \rightarrow N/4 \rightarrow N/8 \rightarrow ... \rightarrow 1$$

- Question: How many times can you divide a number by 2 until you reach 1?
- Answer: log<sub>2</sub>N

$$2 \rightarrow 1$$

$$4 \rightarrow 2 \rightarrow 1$$

$$8 \rightarrow 4 \rightarrow 2 \rightarrow 1$$

$$16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1$$

$$16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1$$

$$32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1$$

$$64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1$$

$$128 \rightarrow 64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1$$

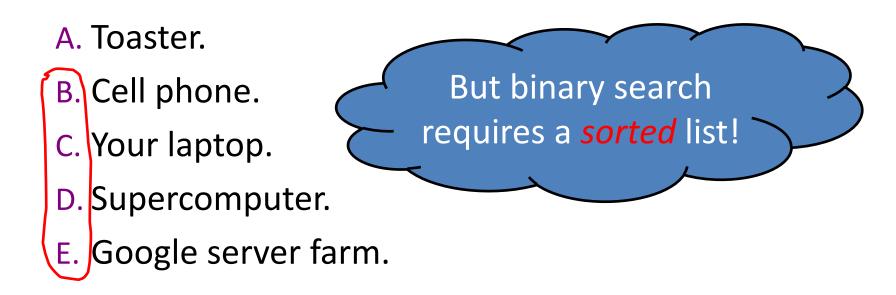
$$256 \rightarrow 128 \rightarrow 64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1$$

$$512 \rightarrow 256 \rightarrow 128 \rightarrow 64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1$$

$$1024 \rightarrow 512 \rightarrow 256 \rightarrow 128 \rightarrow 64 \rightarrow 32 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1$$

# Searching challenge 2

- Problem: A credit card company needs to whitelist 10 million customer account numbers, processing 10,000 transactions per second
- Question: Using <u>binary search</u>, what kind of computer is needed?

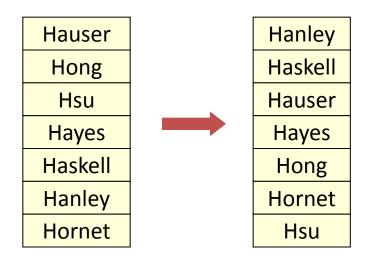


# Sorting

• Sorting problem

Rearrange N items in ascending order

- Applications
  - Statistics, databases, data compression, bioinformatics, computer graphics, scientific computing, ...



### Insertion sort

#### • Insertion sort

- Brute-force sorting solution
- Move left-to-right through array
- Exchange next element with larger elements to its left, one-by-one

2		a							
1	J	0	1	2	3	4	5	6	7
6	6	and	had	him	his	was	you	the	but
6	5	and	had	him	his	was	the	you	but
6	4	and	had	him	his	the	was	you	but
		and	had	him	his	the	was	you	but

Inserting a[6] into position by exchanging with larger entries to its left

#### Insertion sort

#### • Insertion sort

- Brute-force sorting solution
- Move left-to-right through array
- Exchange next element with larger elements to its left, one-by-one

i	4				ä	a			
	j	0	1	2	3	4	5	6	7
		was	had	him	and	you	his	the	but
1	0	had	was	him	and	you	his	the	but
2	1	had	him	was	and	you	his	the	but
3	0	and	had	him	was	you	his	the	but
4	4	and	had	him	was	you	his	the	but
5	3	and	had	him	his	was	you	the	but
6	4	and	had	him	his	the	was	you	but
7	1	and	but	had	him	his	the	was	you
		and	but	had	him	his	the	was	you

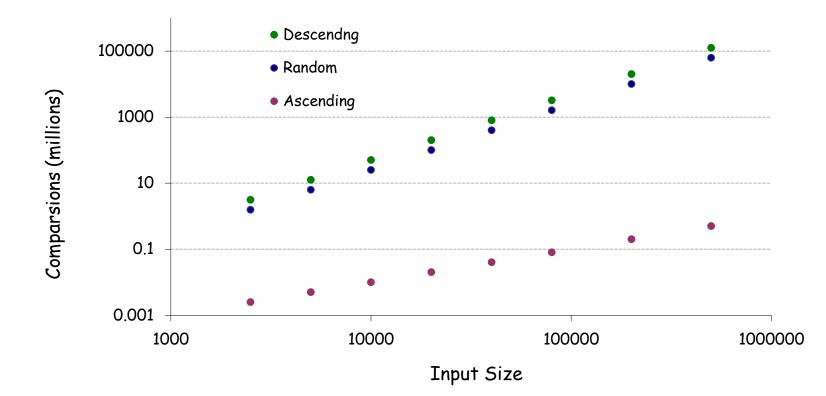
#### Inserting a[1] through a[N-1] into position (insertion sort)

#### Insertion sort: Java implementation

```
public class Insertion
{
    public static void sort(String[] a)
    ł
        for (int i = 1; i < a.length; i++)</pre>
            for (int j = i; j > 0; j--)
                 if (a[j-1].compareTo(a[j]) > 0)
                     exch(a, j-1, j);
                 else break;
    }
    private static void exch(String[] a, int i, int j)
    ł
        String swap = a[i];
        a[i] = a[j];
        a[j] = swap;
    }
```

## Insertion sort: empirical analysis

- Number of compares depends on input family
  - Descending:  $\sim N^2 \, / \, 2$
  - Random:  $\sim N^2 / 4$
  - Ascending:  $\sim N$



## Insertion sort: mathematical analysis

- Worst case [descending]
  - Iteration *i* requires *i* comparisons.
  - Total =  $(0 + 1 + 2 + ... + N-1) \sim N^2 / 2$  compares.



- Average case [random]
  - Iteration i requires i / 2 comparisons on average.
  - Total =  $(0 + 1 + 2 + ... + N-1) / 2 \sim N^2 / 4$  compares



# Sorting challenge 1

- Problem: A credit card company sorts 10 million customer account numbers, for use with binary search.
- Question: Using <u>insertion sort</u>, what kind of computer is needed?
  - A. Toaster.
  - B. Cell phone.
  - C. Your laptop.
  - D. Supercomputer. E. Google server farm.

### Insertion sort: lesson

- Lesson:
  - Even a supercomputer can't rescue a bad algorithm

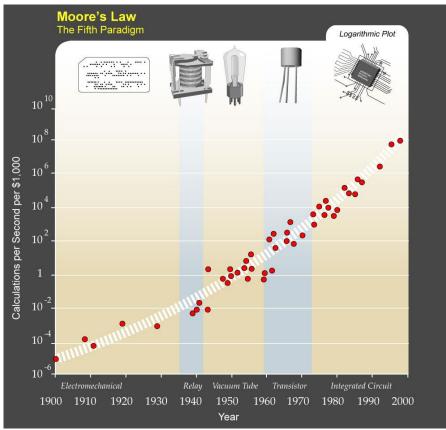
Computer	Comparisons per second	Thousand	Million	Billion
laptop	10 <sup>7</sup>	instant	1 day	3 centuries
super	10 <sup>12</sup>	instant	1 second	2 weeks

## Moore's Law

Moore's law

- Transistor density on a chip doubles every 2 years

- Variants
  - Memory, disk space, bandwidth, computing power per \$



# Moore's law and algorithms

- Quadratic algorithms do not scale with technology
  - New computer may be 10x as fast.
  - But, has 10x as much memory so problem may be 10x bigger
  - With quadratic algorithm, takes 10x as long!

"Software inefficiency can always outpace Moore's Law. Moore's Law isn't a match for our bad coding." – Jaron Lanier



#### Lesson

 Need linear (or linearithmic) algorithm to keep pace with Moore's law

# Mergesort

- Mergesort algorithm
  - Divide array into two halves
  - Recursively sort each half
  - Merge two halves to make sorted whole

input was had him and you his the but sort left and had him was you his the but sort right and had him was but his the you merge and but had him his the was you

# Merging

#### • Merging

Combine two pre-sorted lists into a sorted whole.

- How to merge efficiently?
  - Use an auxiliary array

i	4	k	aux [k]	a							
	J	k	aux[k]	0	1	2	3	4	5	6	7
				and	had	him	was	but	his	the	you
0	4	0	and	and	had	him	was	but	his	the	you
1	4	1	but	and	had	him	was	but	his	the	you
1	5	2	had	and	had	him	was	but	his	the	you
2	5	3	him	and	had	him	was	but	his	the	you
3	5	4	his	and	had	him	was	but	his	the	you
3	6	5	the	and	had	him	was	but	his	the	you
3	6	6	was	and	had	him	was	but	his	the	you
4	7	7	you	and	had	him	was	but	his	the	you

Trace of the merge of the sorted left half with the sorted right half

# Merging

#### • Merging

Combine two pre-sorted lists into a sorted whole.

#### • How to merge efficiently?

Use an auxiliary array

```
String[] aux = new String[N];
// merge into auxiliary array
int i = 10;
int j = mid;
for (int k = 0; k < N; k++)</pre>
{
      (i == mid) aux[k] = a[j++];
   if
   else if (j == hi) aux[k] = a[i++];
   else if (a[j].compareTo(a[i]) < 0) aux[k] = a[j++];</pre>
   else
                                        aux[k] = a[i++];
}
// copy back
for (int k = 0; k < N; k++)
   a[lo + k] = aux[k];
```

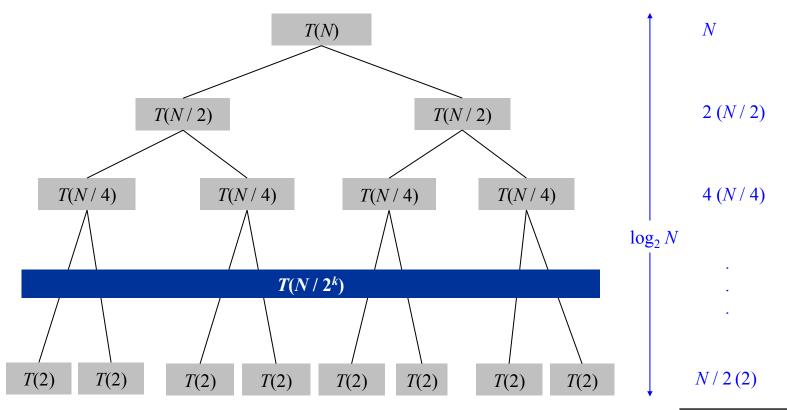
### Mergesort: Java implementation

```
public class Merge
{
    public static void sort(String[] a)
        sort(a, 0, a.length);
    }
    // Sort a[lo, hi).
    public static void sort(String[] a, int lo, int hi)
        int N = hi - lo;
        if (N <= 1) return;</pre>
        // recursively sort left and right halves
        int mid = 10 + N/2;
        sort(a, lo, mid);
        sort(a, mid, hi);
        String[] aux = new String[N];
        // merge sorted halves (see previous slide)
    }
}
```

# Mergesort: mathematical analysis

- Analysis
  - To mergesort array of size N, mergesort two subarrays of size N/2, and merge them together using  $\leq N$  compares

Assume *N* is a power of 2



## Mergesort: mathematical analysis

• Mathematical analysis

analysis	comparisons
worst	$N \log_2 N$
average	$N \log_2 N$
best	$1/2 N \log_2 N$

N	actual	predicted
10,000	120 thousand	133 thousand
20 million	460 million	485 million
50 million	1,216 million	1,279 million

• Validation, theory agrees with observations

# Sorting challenge 2

- Problem: A credit card company sorts 10 million customer account numbers, for use with binary search.
- Question: Using <u>mergesort</u>, what kind of computer is needed?

A. Toaster.

B. Cell phone.

C. Your laptop.

D. Supercomputer.

E Google server farm.

# Sorting challenge 3

 Question: What's the fastest way to sort 1 million 32-bit integers?



http://www.youtube.com/watch?v=k4RRi\_ntQc8

# Mergesort: lesson

- Lesson
  - Great algorithms can be more powerful than supercomputers
  - How long to sort 1 billion things?

Computer	Compares per second	Insertion	Mergesort
laptop	107	3 centuries	3 hours
super	10 <sup>12</sup>	2 weeks	instant

N = 1 billion

# Summary

• Binary search

- Efficient algorithm to search a sorted array

- Mergesort
  - Efficient algorithm to sort an array
- Applications
  - Many many applications are enabled by fast sorting and searching